## 2. OPTIMUM RECEIVER PRINCIPLES

#### 2.1 Maximum Aposteriori Receiver

Consider the generic block diagram of end-to-end communication over the ubiquitous additive white Gaussian noise (AWGN) channel.



• Source:  $\{m_i\}$  with *apriori* probabilities:  $\{P(m_i)\}$ 

• Transmitter: A particular message symbol is represented by a signal waveform allowable in the signal space permitted for a given modulation technique.

$$m = m_i \Leftrightarrow s(t) = s_i(t) \tag{2.1}$$

• Channel: 
$$r(t) = s(t) + n_w(t)$$
 (2.2)

**<u>Problem 1</u>**: Design an optimum receiver which estimates  $\hat{m}$  for the transmitted signal s(t) of a source output *m* such that the probability of error  $P(\varepsilon) \equiv \Pr{ob(\hat{m} \neq m)}$  is MINIMUM.

**Problem 2:** Given that  $\{P(m_i)\}$  are UNKNOWN, which is the real-life problem in many emerging communication systems, design a similar optimum receiver. (Inherently more difficult task!).

**VECTOR CHANNEL:** Consider the case when a sequence of source outputs are bundled into a vector form and transmitted as the case of QAM and other m-ary signaling schemes. In some cases, the signal itself may be in a vector form to start with, as in the case of LPC coefficients in a CELP Speech coder.



• Source Information is mapped into source vectors:  $\{\underline{s}_i; i = 0, 1, ..., M - 1\}$ , where each vector is composed of N-components:  $\underline{s}_i = [s_{i1}, s_{i2}, ..., s_{iN}]$ .

• Received Information is also mapped into vectors:

$$\underline{r}_{i} = \underline{s}_{i} + \underline{n}_{i} = [s_{i1} + n_{i1}, s_{i2} + n_{i2}, \dots, s_{iN} + n_{iN}]$$
(2.3)

Given that the received vector is a point  $\underline{r} = \underline{\rho}$  in the N-dimensional space with coordinates:  $\underline{\rho} = [\rho_1, \rho_2, ..., \rho_N]$  then the optimum receiver must compute the transmitted vector signal  $\underline{s}_i$  for the message  $m_i$  having a maximum aposteriori probability from its knowledge of the set of In other words:

$$\hat{m} = m_k \quad \text{if} \quad \Pr{ob(m_k \mid \underline{r} = \underline{\rho})} > \Pr{ob(m_i \mid \underline{r} = \underline{\rho})} \quad \text{for all } i \neq k \tag{2.4}$$

which is a nearly impossible challenge to meet in many real-life situations.

#### Do we have an equivalent task?

Consider the correct decision for a given incoming vector:

$$\operatorname{Pr}ob(C \mid \underline{r} = \rho) = \operatorname{Pr}ob(m_k \mid \underline{r} = \rho)$$
(2.5)

and the overall correct decision is simply ensemble of correct decisions:

$$\Pr{ob(C)} = \int_{-\infty}^{\infty} \Pr{ob(C \mid \underline{r} = \underline{\rho})} \cdot P_{\underline{r}}(\underline{\rho}) d\underline{\rho}$$
(2.6)

Since  $P_{\underline{r}}(\underline{\rho}) \ge 0$  we do not need to include it in the maximization process, i.e. only the term  $\Pr{ob(C \mid \underline{r} = \rho)}$  must be maximized. Let us use the Bayes Rule on (2.5)

$$\operatorname{Pr} ob(m_i \mid \underline{r} = \rho) = P(m_i) \cdot \operatorname{Pr} ob_r(\rho \mid m_i) / P_r(\rho)$$
(2.7)

but the statement  $m = m_i$  is equivalent to  $\underline{s} = \underline{s}_i$  which implies:

$$\operatorname{Pr}ob_{\underline{r}}(\underline{\rho} \mid m_i) = \operatorname{Pr}ob_{\underline{r}}(\underline{\rho} \mid \underline{s} = \underline{s}_i)$$
(2.8)

Furthermore, the denominator term is independent of the index i, hence, the maximization and we have the revised principle for our optimum receiver:

$$\hat{m} = m_k$$
 if  $P(m_i)$ . Pr  $ob_{\underline{r}}(\underline{\rho} | \underline{s} = \underline{s}_i)$  is maximum when  $i = k$  (2.9)

When  $P(m_i)$  are not known and the receiver can only maximize the last portion of (2.9). Then we have a restricted version of the general optimum receiver called **MAXIMUM-LIKELIHOOD** (ML) Receiver.

#### **ML Receiver Principle:**

$$\hat{m} \Rightarrow m_k$$
 when  $\operatorname{Pr}ob_{\underline{r}}(\underline{\rho} \mid \underline{s} = \underline{s}_i)$  is MAXIMUM. (2.10)

Decision Regions are needed to perform the mapping properly for each signal vector.

**Example 2.1:** Given (3) input vectors in a 2-D vector space with the following signal set assignment:  $m_0 \Rightarrow \underline{s}_0 = [1,2]; \quad m_1 \Rightarrow \underline{s}_1 = [2,1];$  and  $m_2 \Rightarrow \underline{s}_2 = [1,-2]$ 



Let us also assume that the input message probabilities:  $P(m_0), P(m_1), P(m_2)$  are given. For this assignment, our receiver will compute:

#### An ML Receiver will choose the index of the message with the largest product above.

For every point  $\underline{\rho}$  in  $(\varphi_1, \varphi_2)$  plane an assignment can be made if the plane is partitioned into disjoint regions  $\{I_i\}$  for i=0,1,2; which are called decision regions, very similar to the codeword selection process in Vector Quantization (VQ). Then we have the ML receiver as a simple geometric map:

$$\underline{r} = I_k \implies \hat{m} = m_k$$
 and an error is made if  $\hat{m} \Rightarrow m_k$  iff  $\underline{r} \not\subset I_k$  (2.11)

# 2.2 ML Receiver for AWGN Channel

Given that the signal in the channel is corrupted by a zero mean AWGN with a variance  $\sigma^2$ .

$$\underline{r} = \underline{s} + \underline{n} = [s_1 + n_1, s_2 + n_2, \dots, s_N + n_N]$$
(2.12)

Now:

$$\underline{r} = \underline{\rho} \text{ when } \underline{s} = \underline{s}_i \quad iff \quad \underline{n} = \underline{\rho} - \underline{s}_i \tag{2.13}$$

And then

$$P_{\underline{r}}(\underline{\rho} | \underline{s} = \underline{s}_i) = P_{\underline{n}}(\underline{\rho} - \underline{s}_i | \underline{s} = \underline{s}_i) \quad \text{for} \quad i = 0, 1, \dots, M - 1$$

$$(2.14)$$

Since the signal <u>s</u> and the channel noise <u>n</u> are statistically independent  $P_{\underline{n}|\underline{s}} = P_{\underline{n}}$ . This simplifies (2.14) into:

$$P_{\underline{n}}(\underline{\rho} - \underline{s}_i \mid \underline{s} = \underline{s}_i) = P_{\underline{n}}(\underline{\rho} - \underline{s}_i)$$
(2.15)

In this case, the general ML decision function becomes  $P(m_i) \cdot P_n(\underline{\rho} - \underline{s}_i)$ . Now the components of signal is assumed to be independent, noise has a zero-mean we can write the noise distribution:

$$P_{\underline{n}}(\underline{u}) = \frac{1}{(2\pi\sigma^2)^{N/2}} \exp\{-\frac{1}{2\sigma^2} \sum_{j=1}^N u_j^2\}$$
(2.16)

Let us use the following dot-product notation:

$$\left|\underline{u}\right|^2 = \underline{u} \bullet \underline{u}^* = \sum_{j=1}^N u_j^2$$

Our distribution is written as:

$$P_{\underline{n}}(\underline{u}) = \frac{1}{(2\pi\sigma^2)^{N/2}} \exp\{-\frac{1}{2\sigma^2} |\underline{u}|^2\}$$
(2.17)

Then for this probability system we have the ML principle as:

$$\hat{m} \Rightarrow m_i \quad whenever \quad P(m_i).\exp\{-\frac{1}{2\rho^2} \left|\underline{\rho} - \underline{s}_i\right|^2\} \text{ is maximum.}$$
 (2.18)

Equivalently, the task is to **MINIMIZE**:

$$\left|\underline{\rho} - \underline{s}_i\right|^2 - (2\sigma^2) \cdot \log_e P(m_i) \tag{2.19}$$

The first term is the Euclidean Distance between the received vector and a candidate signal vector. If all the messages are equally likely then the optimum decision rule does not depend on the index at all and we have the **MINIMUM MEAN-SQUARE (MMS) DISTANCE Receiver**. That is we

assign the message index of the closest neighbor of the incoming vector, which is also known as the Nearest Neighbor Rule in VQ and other clustering techniques.

# 2.3 Correlation and Matched-Filter Receivers

If we revisit the Communication System Block Diagram for vector signals as shown below, it would be necessary to synthesize waveform signals to be transmitted over real-life channels, such as the twisted-pair or coaxial cable, microwave or fiber-optic links.



- It is necessary to synthesize the signal set  $\{s_i(t)\}$  at the transmitter. This can be achieved by "building blocks waveforms".
- Synthesis signal sets and Recovery of signal vectors:
- 1. A set of N integrating filters are used to generate N signal components with strengths  $\{s_{ii}\}$ .
- 2. The filter outputs are summed to yield the signal waveform: s(t) to be transmitted for a particular message  $m_i$  for each of M different messages.



1. Let us choose the building-block waveforms from an orthonormal set such that:

$$\int_{-\infty}^{\infty} \varphi_j(t) \varphi_l(t) dt = \begin{cases} 1 & \text{if } j = l \\ 0 & \text{if } j \neq l \end{cases} \quad \text{for all} \quad 1 \le i, j \le N$$

$$(2.21)$$

- 2. This will yield a probability of error independent of the actual wave-shapes.
- 3. We can exactly recover the signal vectors and hence, the messages in the absence of channel if we push these synthesized waveforms of (2.20) into a simple integrating filter structure as shown above.

$$\int s_{i}(t)\varphi_{l}(t)dt = \int \left[\sum_{j=1}^{N} s_{ij}\varphi_{j}(t)\right]\varphi_{l}(t)dt = \sum_{j=1}^{N} s_{ij}\delta_{jl} = s_{il}$$
(2.22)

If we perform similar integration for all the branches we obtain:  $s_i = [s_{i1}, s_{i2}, ..., s_{iN}]$ .

- 4. Examples of Orthonormal Signal Sets:
  - Orthonormal time-shifted pulses:  $\varphi_j(t) = g(t j\tau)$  for j = 1, 2, ..., N
  - Orthonormal Fourier Transform pulses:  $\varphi_j(t) = \begin{cases} \sqrt{1/\tau} & -T \le \tau < 0\\ 0 & otherwise \end{cases}$





The optimum ML receiver of the system performs:

Set:  $\hat{m} = m_k \quad if \left| \underline{r} - \underline{s}_i \right|^2 - 2\sigma_n^2 \log P(m_i)$  is MINIMUM. (2.23)

Square operations can be eliminated in (2.23) by observing:

$$\left|\underline{r} - \underline{s}_{i}\right|^{2} = \left|\underline{r}\right|^{2} - 2(\underline{r} \bullet \underline{s}_{i}) + \left|\underline{s}_{i}\right|^{2}$$

$$(2.24)$$

where the dot product is given also by:

$$\underline{r} \bullet \underline{s}_{i} \equiv \sum_{j=1}^{N} r_{j} \underline{s}_{ij}$$
(2.25)

## **Observations**:

• Note 1: First term in (2.24) is independent of the index and no need to worry in optimization.

• Note 2: Last terms in (2.23) and (2.24) depend only on source side information supplied by the designer then they can be combined into a constant parameter set and burned into the ROM of the system:

$$c_{i} = (1/2)[\sigma_{n}^{2}\log P(m_{i}) - |\underline{s}_{i}|^{2}]$$
(2.26)

• The optimum receiver of (2.23) is now equivalent to:

Set: 
$$\hat{m} = m_k$$
 if  $(\underline{r} \bullet \underline{s}_i + c_i)$  is MAXIMUM. (2.27)

which is simply the structure of a **CORRELATION RECEIVER**.



**Note:** When the source vocabulary size M is not very large this implementation is not costly and most of the operations to the right of the *"Integrators*" can be done by table look-ups. However, when M is very large then the dot-products are usually handled by using "DSP" based devices. The use of multipliers can be avoided if we replace the structure to the left of the *"Weighting Matrix"* as follows:

1. Let us consider a filter with an impulse response

$$h_{j}(t) = \varphi_{j}(T-t). \qquad (2.28)$$

$$+ \sqrt{\frac{2}{T}} \int_{0}^{\varphi_{j}(t)} h_{j}(t) = \varphi_{j}(T-t)$$

$$- \sqrt{\frac{2}{T}}$$

2. If the input to this filter is r(t) then its response is simply:

$$u_{j}(t) = \int_{-\infty}^{\infty} \underline{r}(\alpha)h(t-\alpha)d\alpha = \int_{-\infty}^{\infty} \underline{r}(\alpha)\varphi_{j}(T-t+\alpha)d\alpha$$
(2.29)

3. When we sample the output at t=T we have

$$u_j(T) \equiv r_j \tag{2.30}$$

4. Finally, the task is to push it through the weighting matrix and the rest of the receiver above.

This receiver is called a "MATCHED-FILTER" Receiver since it is constructed by using the shifted versions of the signal building block functions:  $\varphi_i(T-t)$ .



**Example 2.2:** Consider the case for a gated-sinusoidal tone signal with a gate period of T seconds as shown below. The convolution operation in the matched -filter above will result in a triangular enveloped sinusoid with the same frequency and thus it will peak at the sampling instant T.

